

# Max Intersection-Complete Codes

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# Motivation

- The 2014 Nobel Prize in Physiology or Medicine was awarded for the discovery of place cells and grid cells
- Place cells represent an animal's location
- Multiple place cells can fire at once

## Definition

A **neural code**  $\mathcal{C}$  on  $n$  neurons is a set of subsets of  $[n]$ .

- Given  $n$  neurons, we build neural codes from their respective *receptive fields*, living in  $\mathbb{R}^d$ .
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- The receptive field of a neuron  $i$  is denoted  $U_i$ .
- On 5 neurons, one codeword could be  $\{2,4\}$ ; this is where the receptive fields  $U_2$  and  $U_4$  overlap; we write this as 24.

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- Certain types of codes are known to be convex, notably max intersection-complete codes.

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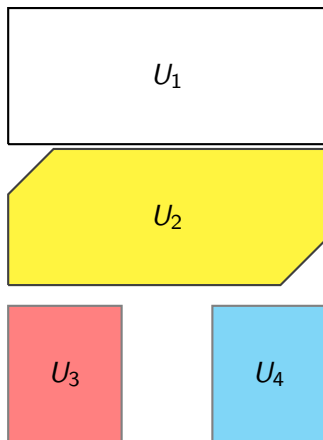
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The **maximal code** on  $n$  neurons is  $\mathcal{C}_{max}(n) = \{\sigma : \sigma \subseteq [n]\}$ .

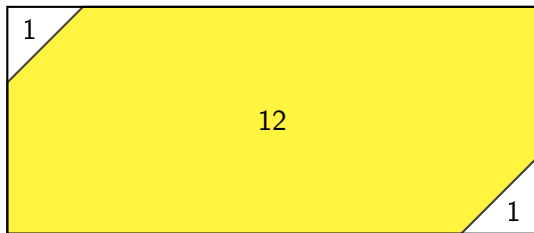
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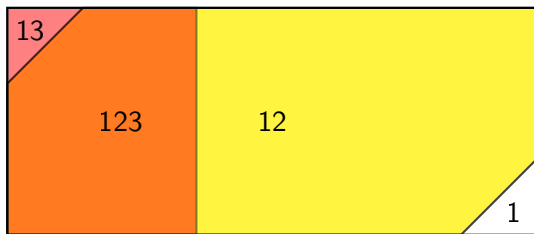
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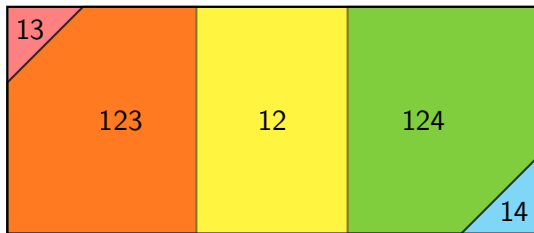
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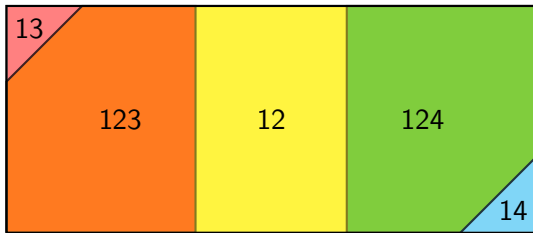


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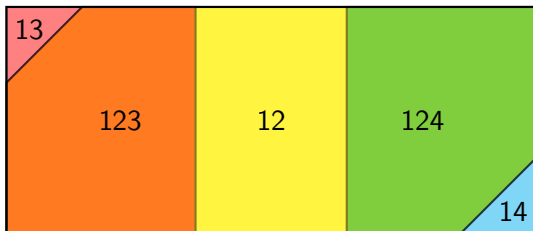
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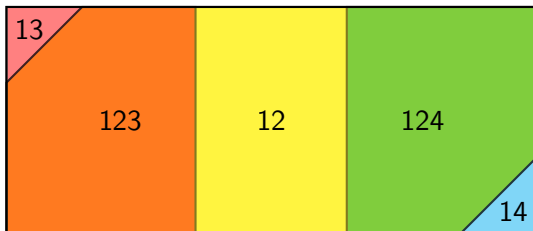


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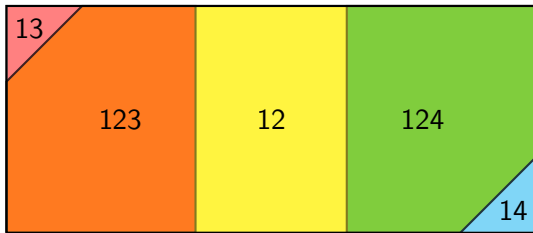




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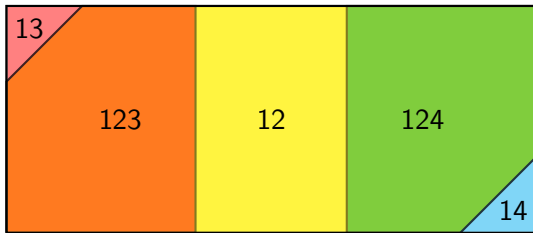
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Max intersection-complete? Yes!

From a neural code  $\mathcal{C}$ , we obtain its neural ideal  $J_{\mathcal{C}}$ , defined to be

$$J_{\mathcal{C}} := \left\langle \prod_{i \in \sigma} x_i \prod_{j \in \tau} (1 - x_j) : \sigma \notin \mathcal{C}, \tau = [n] - \sigma \right\rangle.$$

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In our example, 24 is not a codeword of  $\mathcal{C}$ , so

$$x_2 x_4 (1 + x_1)(1 + x_3) \in J_{\mathcal{C}}$$

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- Type 1 relations:  $\prod_i x_i$
- Type 2 relations:  $\prod_i x_i \prod_j (1 - x_j)$
- If a Type 1 relation  $x_{a_1} \dots x_{a_n}$  is in the CF of  $J_C$ , then the codeword  $c = a_1 \dots a_n$  is not in  $\mathcal{C}$ , nor is any codeword containing  $c$ .
- If a Type 2 relation  $x_{a_1} \dots x_{a_n} (1 - x_{b_1}) \dots (1 - x_{b_m})$  is in the CF, then

$$\bigcap_{i \in \{a_1, \dots, a_n\}} U_i \subseteq \bigcup_{j \in \{b_1, \dots, b_m\}} U_j$$

# Canonical Form Example

Recall our code  $\mathcal{C} = \{123, 124, 12, 14, 13, \emptyset\}$ .

Here,  
 $CF(\mathcal{C}) = \{x_2(1-x_1), x_3(1-x_1), x_4(1-x_1), x_3x_4, x_1(1-x_2)(1-x_3)(1-x_4)\}$ .

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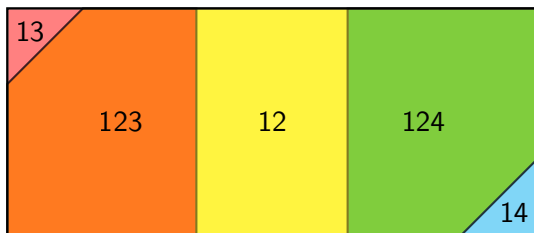
Further, an element like  $x_2(1-x_1)$  tells us that  $U_2 \subseteq U_1$ .

# Canonical Form Example

Similarly, because  $x_1(1 - x_2)(1 - x_3)(1 - x_4) \in CF(J_C)$ , we have that  $U_1 \subseteq U_2 \cup U_3 \cup U_4$ .

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# An Existing Signature for Intersection-Complete Codes

The following theorem gives a signature in the canonical form for intersection-complete codes:

## Theorem (Curto, Gross, et al. 2015)

A code  $\mathcal{C}$  is intersection-complete if and only if  $CF(J_{\mathcal{C}})$  contains only monomials and pseudomonomials of the form  $(1 - x_j) \prod_i x_i$ .

## Research Question

Does there exist a signature in the canonical form for maximum intersection-complete codes?



# Finding the Facets

We have been able to develop an algorithm for finding the facets of a code  $\mathcal{C}$  from  $CF(J_{\mathcal{C}})$ .

We use the fact that if a monomial appears in  $CF(J_{\mathcal{C}})$  then no codeword containing the indices of that monomial appears in  $\mathcal{C}$ .

# Example of Facet Algorithm

Recall our earlier example:  $\mathcal{C} = \{\mathbf{123}, \mathbf{124}, 12, 13, 14, \emptyset\}$ .

The only monomial in  $CF(J_{\mathcal{C}})$  is  $x_3x_4$ .

On 4 neurons,

$\mathcal{C}_{max} = \{\emptyset, 1, 2, 3, 4, 12, 13, 14, 23, 24, 34, 123, 124, 134, 234, 1234\}$ .

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Removing all codewords eliminated by this monomial gives us

$\mathcal{C}'_{max} = \{\emptyset, 1, 2, 3, 4, 12, 13, 14, 23, 24, \mathbf{34}, 123, 124, \mathbf{134}, \mathbf{234}, \mathbf{1234}\}$ .

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$$\mathcal{C}'_{max} = \{\emptyset, 1, 2, 3, 4, 12, 13, 14, 23, 24, \mathbf{123}, \mathbf{124}\}.$$

We see that the facets of  $\mathcal{C}$  and our reduced  $\mathcal{C}'_{max}$  are the same.

# Sufficient Condition for Non-maximality

## Proposition (Franke-H)

Let  $\mathcal{C}$  be a neural code,  $J_{\mathcal{C}}$  be its neural ideal, and  $CF(J_{\mathcal{C}})$  be the corresponding canonical form. If there exist  $\tau \subset [n]$  and  $\sigma \subseteq [n] - \tau$  such that  $\prod_{i \in \tau} x_i \in CF(J_{\mathcal{C}})$  and  $\prod_{j \in \sigma} x_j \prod_{i \in \tau} (1 - x_i) \in CF(J_{\mathcal{C}})$ , then  $\mathcal{C}$  is not convex.

## Corollary

For a code to be max intersection-complete, it cannot have the above condition.

# Example

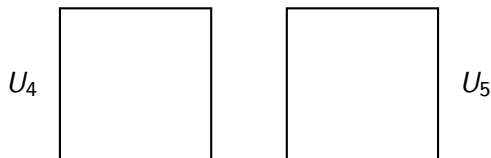
Let  $\mathcal{C} = \{4, 5, 1234, 1235, \emptyset\}$  be a code on five neurons. The canonical form contains both  $x_4x_5$  and  $x_1(1 - x_4)(1 - x_5)$ .

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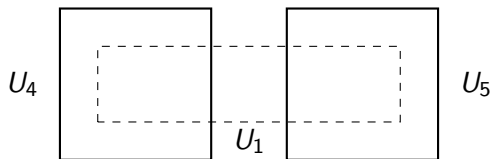
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# An Algorithm for Determining Missing Codewords

1. Pick a complex pseudomonial

This is a pseudomonial with multiple  $(1 - x_j)$  factors, e.g.,  
 $x_i(1 - x_{j_1}) \dots (1 - x_{j_m})$ . Write  $\bigcap_{k \in [m]} ij_k = i$ .

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## 2. Add “equivalent” neurons

Neurons which always fire together are equivalent, e.g. if  $x_i(1 - x_j)$  and  $x_j(1 - x_i) \in CF(J_C)$ , then neurons  $i$  and  $j$  are equivalent.

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## 3. Add all other possible neurons not prevented by monomials.

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2.  $x_4(1 - x_5), x_5(1 - x_4), x_2(1 - x_3), x_3(1 - x_2) \in CF(J_{\mathcal{C}})$ , so

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3.  $x_4x_6, x_5x_6$  are the only monomials in  $CF(J_{\mathcal{C}})$ , so we can add 1 to 2345  
and 236 to get  $1236 \cap 12345 = 123$

# Acknowledgments

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